



FraDeCo(P)P 2012

PROOFS OF POINTER ALGORITHMS BY INDUCTIVE
REPRESENTATION OF GRAPHS

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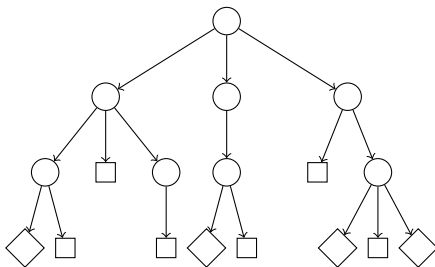
Outline

- 1 Context
 - Introduction
 - Approach
 - Verification of graph transformations
- 2 Schorr-Waite
 - Demo
 - Description
 - Verification
- 3 BDD
 - Description
 - Verification
- 4 Conclusion

First remark

Most usual data structures are tree-shaped with several kind of additional pointers:

- sharing pointers ->
- root pointers ->
- father pointers ->
- ...

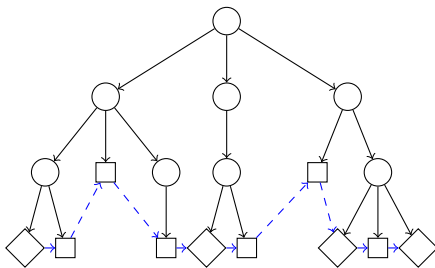


For numerous algorithms, different views on data-structures can be used for the verification

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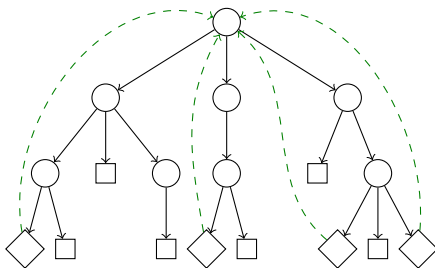


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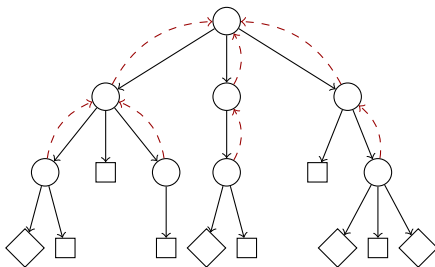


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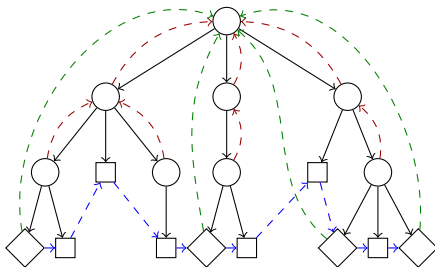


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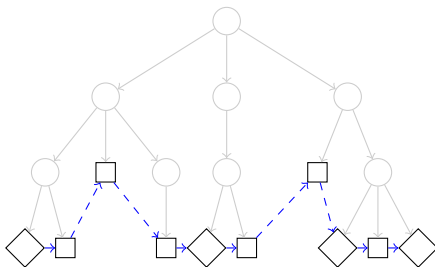


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For numerous algorithms, different views on data-structures can be used for the verification

Verification

Objective

Generate efficient and verified programs
(with pointers and/or mutable structures)

Safe code generation

```
graph TD; Objective[Objective] --> Safe[Safe code generation]; Safe --> Abstract[Abstract Language]; Safe --> Imperative[Imperative language];
```

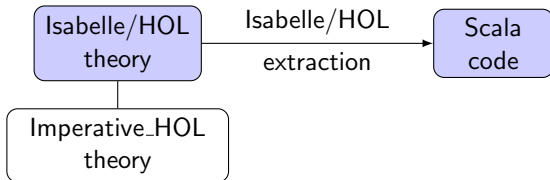
Abstract Language

- Arborescent data structure
+ pointers
- Properties verified by induction

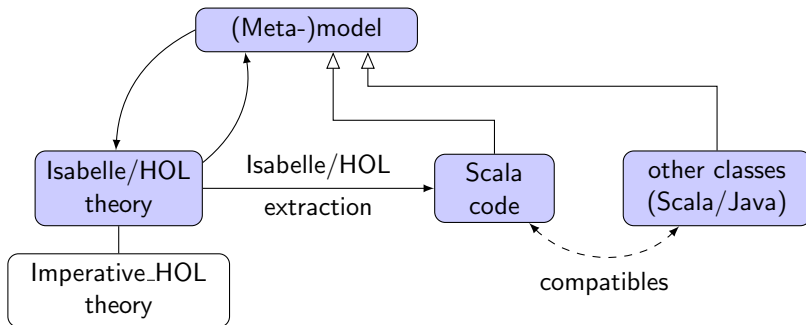
Imperative language

- Efficient :
 - Sharing
 - Mutability
- Same properties

General frame



General frame



Related work

Graph representation

- Nodes and edges
- Coinduction
- Trees + pointers:
 - PALE [MS01]
 - Term-graph rewriting in Tom using relative positions [BB08]
 - Locally nameless encoding of lambda-terms [Cha09]
 - B+ trees functional representation with pointers [MM10]

Verification

- Proofs on raw heap and pointers
- Hoare (Separation) Logic [Rey02]

Case studies - Overview

Schorr-Waite

- Arbitrary rooted graph (with outgoing arity ≤ 2)
- High mutability
- Proof by simulation
- LOPSTR'2010 [GSMP10]

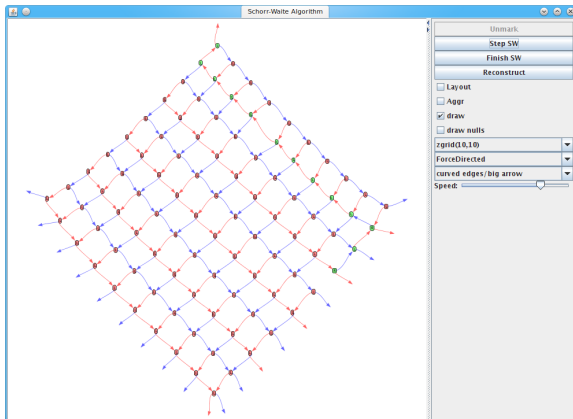
BDD

- Rooted acyclic graph (shared tree)
- References comparisons
- Direct proof
- TAAPSD'2010 [GS10], FoVeOOS'2011 [GS11]

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Demo



Algorithm features

Purpose

- Marking graphs without using more space (stack, ...)
- Traversing a tree by terminal recursivity and without stack

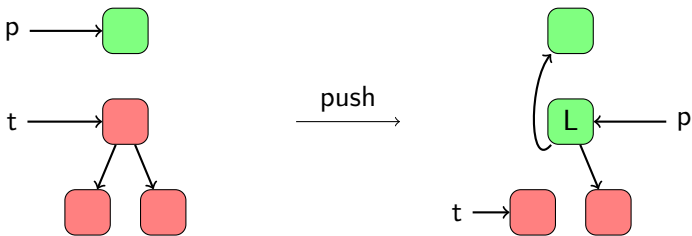
Use

Garbage collector, case study...

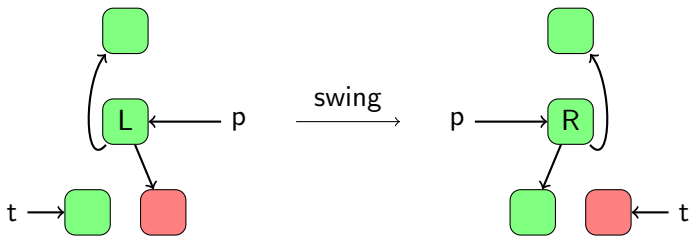
Principle

- Modification of the graph pointers to store the path to the root
- 2 variables containing the pointers :
 - t : to the current node
 - p : to the previously visited node

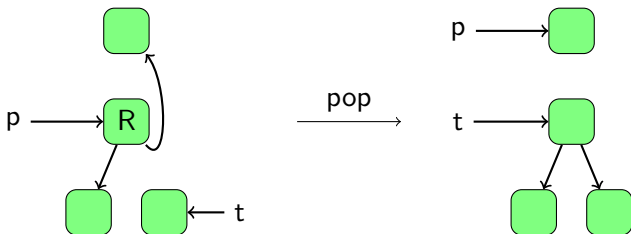
Steps : Push



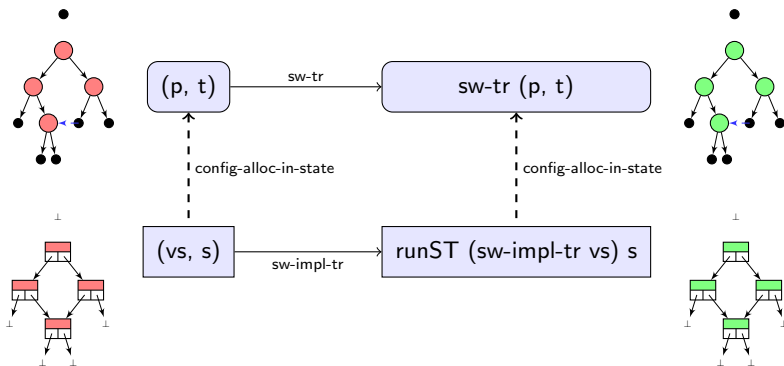
Steps : Swing



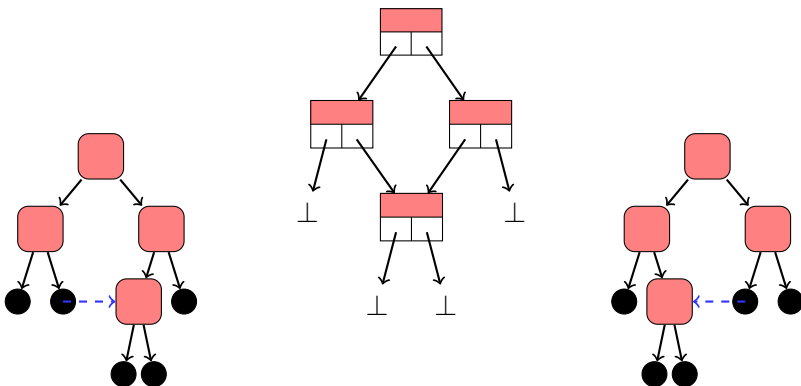
Steps : Pop



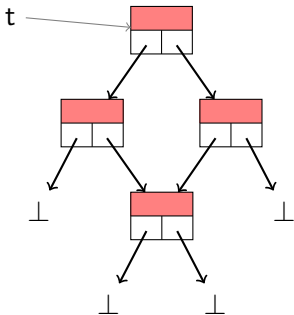
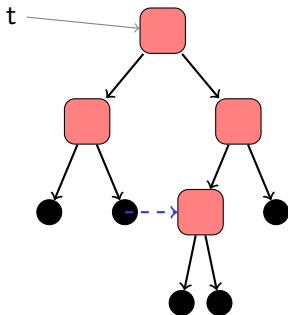
Proof by simulation



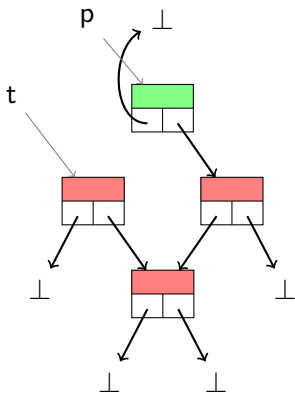
Choosing the spanning tree: Example with 2 possibility



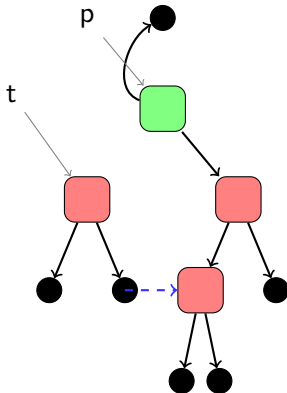
The wrong one!

 $p \rightarrow \perp$  $p \rightarrow \bullet$ 

The wrong one!

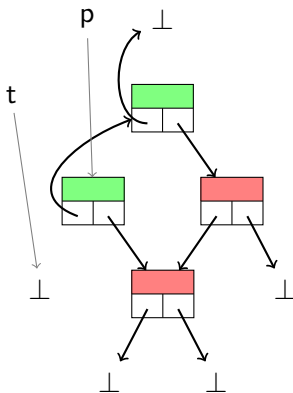


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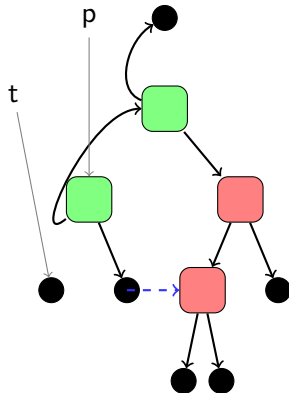


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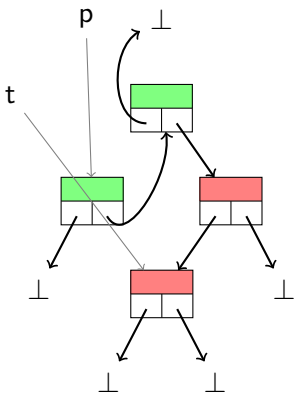


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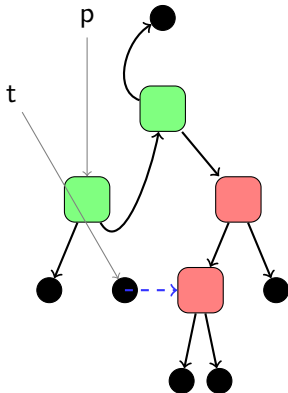


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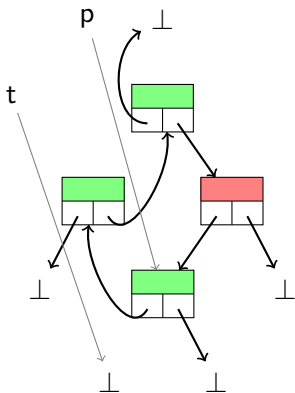


swing

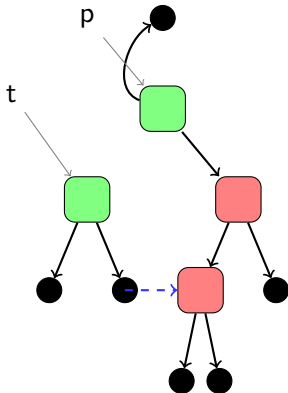


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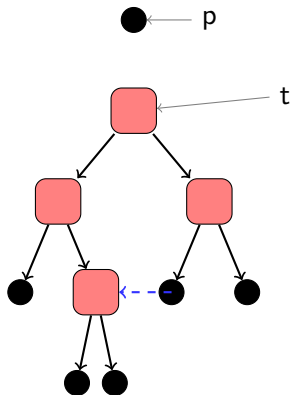
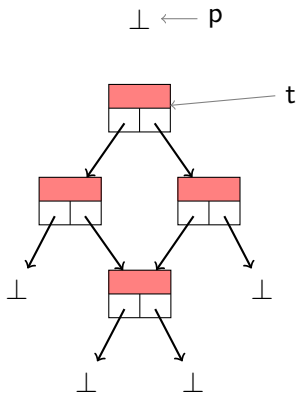


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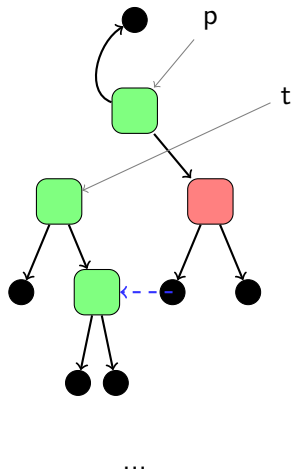
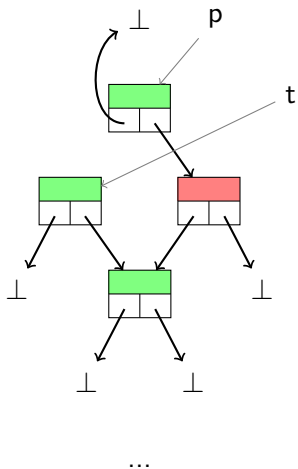


pop

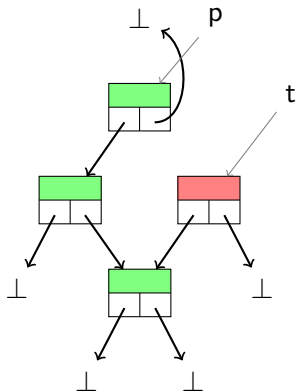
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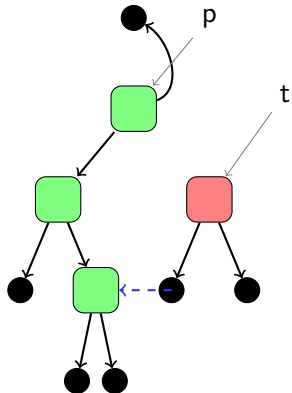
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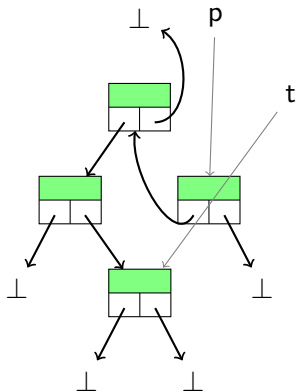


swing

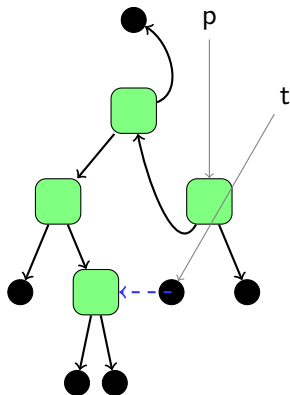


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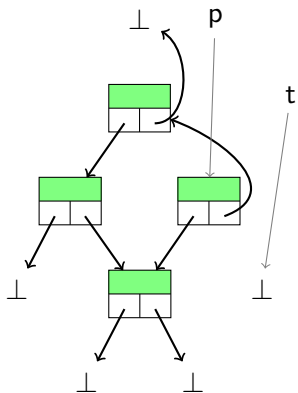


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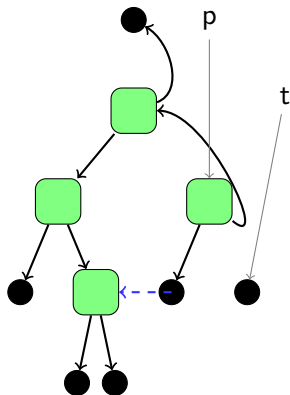


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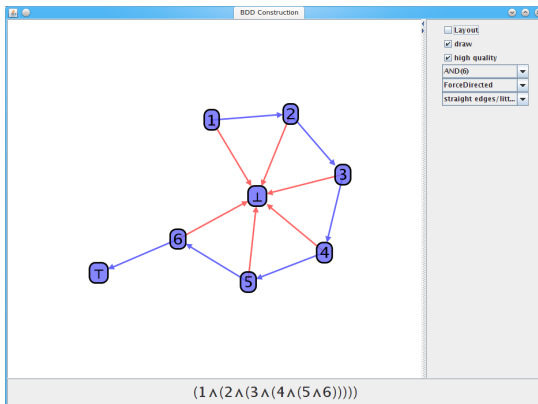


swing

Outline

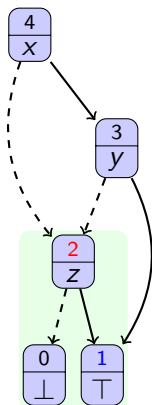
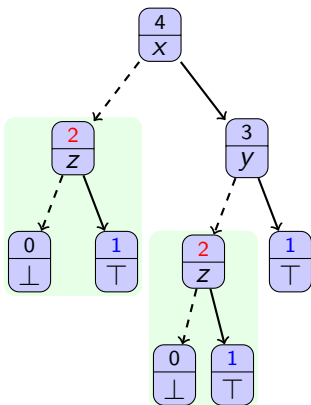
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Demo



Sharing

We add references to represent sharing. They also allow to add mutable content in the nodes.



0	nbrefs = 0 ...
1	nbrefs = 0 ...
2	nbrefs = 0 ...
3	nbrefs = 0 ...
4	nbrefs = 1 ...
5	nbrefs = - ...

⋮

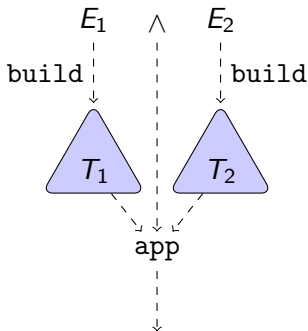
build

primrec build ::

$\text{'v expr} \Rightarrow (\text{bool}, \text{'v}) \text{ rtree Heap}$

where

```
build (Var i) = do{
  cf ← constLeaf False;
  ct ← constLeaf True;
  mk i cf ct
}
| build (Const b) = (constLeaf b)
| build (BExpr bop e1 e2) = do{
  n1 ← build e1;
  n2 ← build e2;
  app bop (n1, n2)
}
```

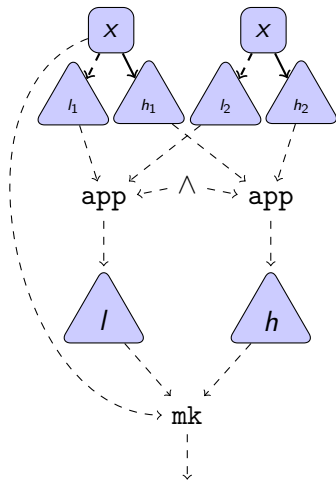


app

```

function app ::
  (bool  $\Rightarrow$  bool  $\Rightarrow$  bool)
 $\Rightarrow$  ((bool, 'v) rtree * (bool, 'v) rtree)
 $\Rightarrow$  (bool, 'v) rtree Heap
where
  app bop (n1, n2) = do {
    if tpair is-leaf (n1, n2) then
      constLeaf (bop (leaf-val n1) (leaf-val n2))
    else do {
      let ((l1, h1), (l2, h2)) =
          select split-lh dup (n1, n2);
          l  $\leftarrow$  app bop (l1, l2);
          h  $\leftarrow$  app bop (h1, h2);
          mk (varOfLev (min-level (n1, n2))) l h
      }
    }
  }

```



Memoization and Garbage Collection

Memoization

Records previous computations results to reuse them

- little change in functions and proofs

Garbage collection

Removes no more used BDDs from the maximal sharing table

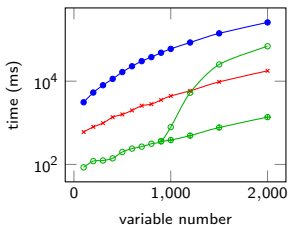
- several reference counter mutations in functions and proofs
- weakening of an invariant

Main theorem & benchmarks

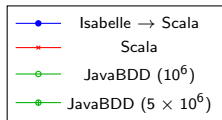
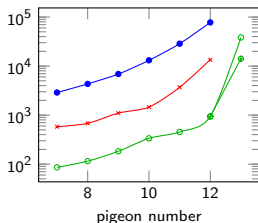
Equivalent expressions construct the same BDD

lemma build-correct: $\llbracket \forall t \in \text{trees } s1 \cup \text{trees } s2. \text{robdd-refs } t;$
 $\text{wf-heap } s1; \text{ effect } (\text{build } e1) s1 s1' t1;$
 $\text{wf-heap } s2; \text{ effect } (\text{build } e2) s2 s2' t2 \rrbracket$
 $\implies (\text{interp-expr } e1 = \text{interp-expr } e2) \longleftrightarrow \text{struct-equal } (t1, t2)$

Urquhart benchmark



Pigeonhole benchmark



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Approach

Representing graphs as (trees + pointers) to verify transformations

Schorr-Waite

- Arbitrary rooted graph (with outgoing arity ≤ 2)
- High mutability / Proof by simulation

BDD

- Rooted acyclic graph (shared tree)
- References comparisons / Direct proof

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Thanks for your attention



Emilie Balland and Paul Brauner.

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